### Week 13-15: Engineering II – Shannon's Information Theory

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# Claude E. Shannon (1916-2001)<sup>\*</sup>: Father of Information Theory

- An American mathematician, electrical engineer, and cryptographer.
- Founded information theory with a landmark paper "<u>A Mathematical Theory of</u> <u>Communication</u>" published in 1948.
- Provided a framework where we can study how to handle information, i.e.,
  - Quantification
  - Storage
  - Communication



## Mathematicians Active in Science & Technology



#### Shannon's Mouse





#### Fundamental Questions

- What is the ultimate data compressions?
  - Answer: Entropy H.
- What is the ultimate transmission rate of communication?
  - Answer: Channel capacity **C**.

#### Information and Entropy

• The entropy of H(X) of a discrete random variable X with alphabet X and probability mass function  $p(x) = \Pr\{X = x\}, x \in X$ , is defined as

$$H(X) \stackrel{\text{def}}{=} \sum_{x \in X} \{-\log p(x) \times p(x)\} = E[-\log p(x)]$$

- The entropy is the *measure of uncertainty* of a random variable and interpreted as the expected value of the information carried by each alphabet, i.e.,  $-\log p(x)$  (shown in the figure).
  - Entropy is expressed in **bits** when the log is to the base 2, i.e.,  $-\log_2 p(x)$ .



#### Huffman Encoding

- Entropy of the source: 1.74 bits/symbol
- Average length of binary coding: 2 bits/symbol
- Average length of Huffman coding: 1.85 bits/symbol



Symbol	Probability	Binary Coding	Huffman Coding
김	0.4	00	0
0	0.35	01	10
박	0.2	01	110
최	0.05	11	111

#### Channel Capacity

$$C = B \log_2\left(1 + \frac{S}{N}\right)$$

- C: Maximum channel capacity [bits/second]
- B: Bandwidth of the channel [Hz]
- S: Signal power [Watts]
- *N*: Noise power [Watts]

#### Capacity of Optical Fiber

• It is known that the maximum capacity of one strand of optical fiber is greater than 100 Terabits/second<sup>\*</sup>, i.e.,

 $C > 100 \times 10^{12} bits/second$ 

- The size of 2-hour-long HD-quality movie is around 4 Gigabytes, i.e.,  $32 \times 10^9$  bits.
  - The download time is as follows:

 $\frac{32 \times 10^9}{100 \times 10^{12}} = 32 \times 10^{-5} s = 320 \ \mu s = 0.32 \ ms$ 

\* Mitra & Stark, Nature, vol 411, June 28, 2001.

